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# Analytical Model of IEEE 802.15.4 for Multi-hop Networks

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# Motivations

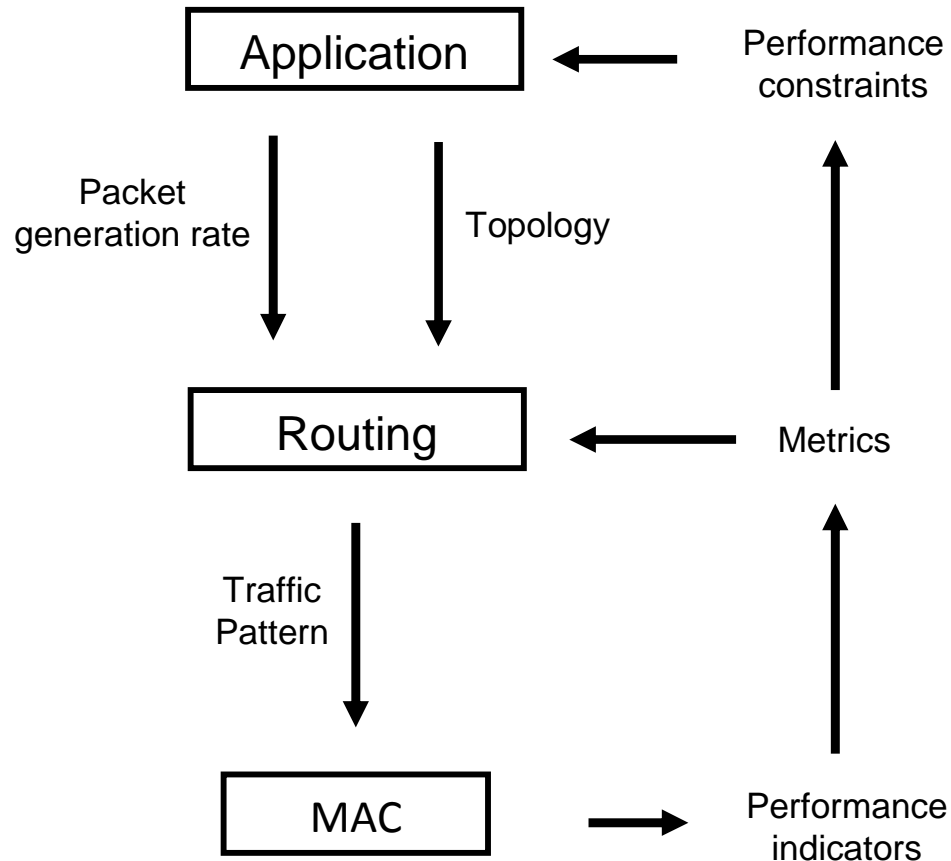
- **Distributed** wireless applications set new requirements and limitations on network protocols
    - High reliability, low latency, energy efficiency
  - **Multi-hop** communication is instrumental for many application scenarios.
  - Ongoing standardization for **routing** over low-power lossy networks (ROLL) by IETF.
  - IEEE 802.15.4 standard for **MAC** and physical layer in WPAN.
- Analytical study of the interaction between routing decision and MAC performances of IEEE 802.15.4 in multi-hop networks with ROLL specifications.





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# Protocol Interactions



- RPL - Routing over low power and lossy networks
  - DODAG (Destination-oriented Directed Acyclic Graph) structure
  - Multipoint-to-point communications
- Unslotted IEEE 802.15.4 MAC
  - Random access based on CSMA/CA with binary exponential backoff (BEB)

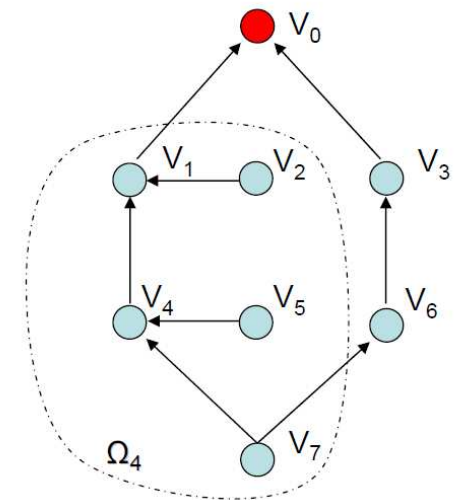




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# System Model

- Key aspects of multihop networks
  - Heterogeneous traffic condition
  - Limited carrier sensing range (hidden terminal problem)
- Topology
  - $\Omega_i$ : neighborhood set of node  $i$
  - $\Omega_{j/i}$ : hidden terminal set of link  $(i,j)$
  - $\Gamma_i$ : parent set of node  $i$
  - $\Delta_i$ : children set of node  $i$

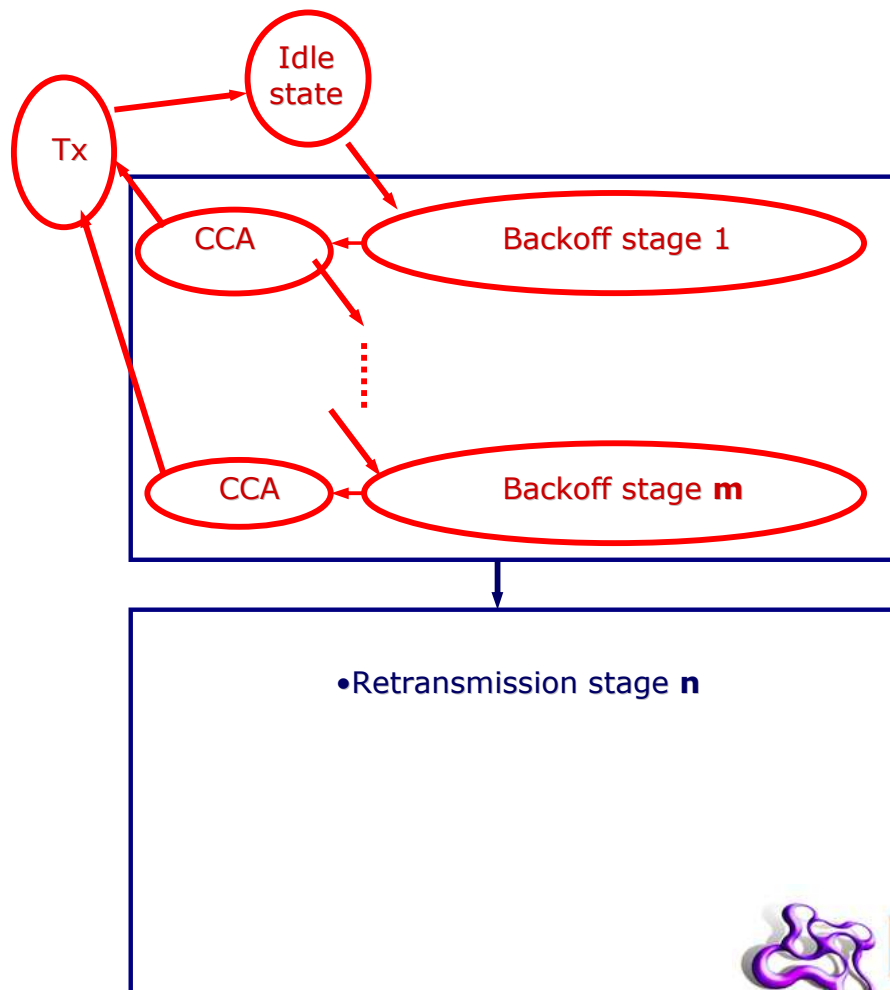




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# Unslotted IEEE 802.15.4 MAC

## • Binary Exponential Backoff (BEB):

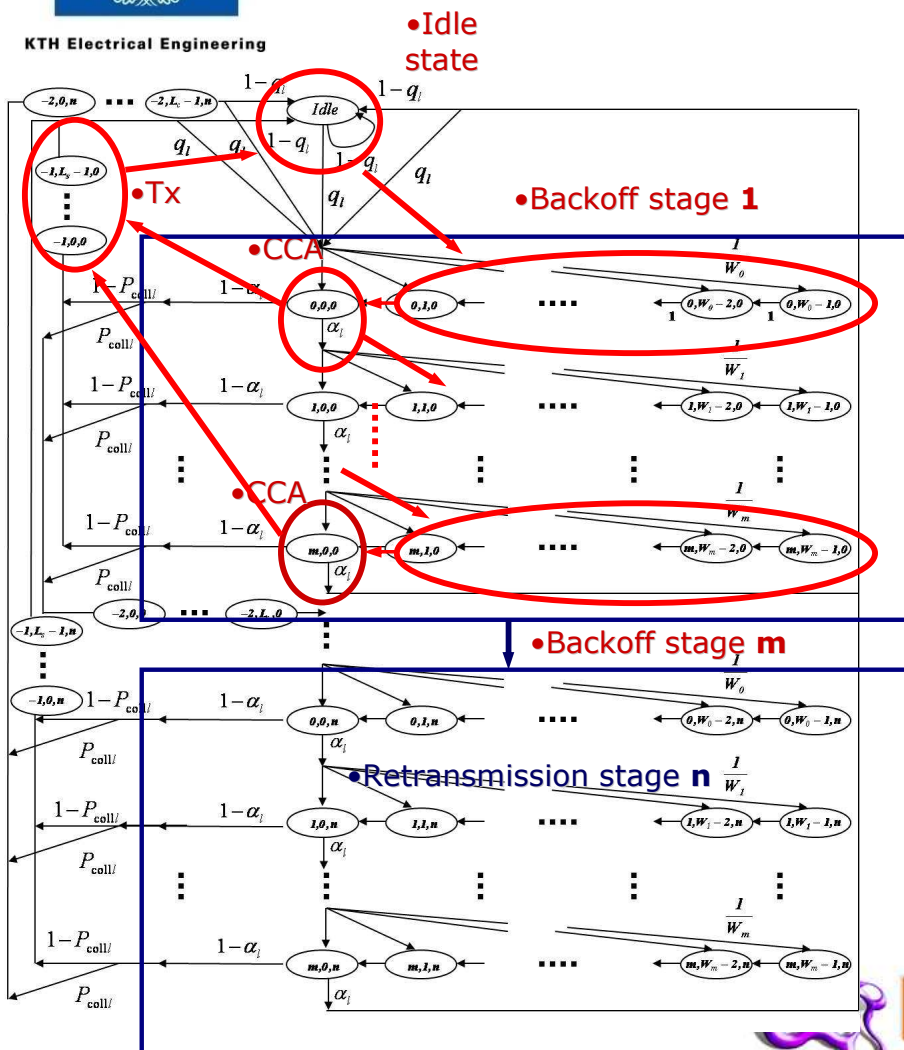


- A transmitting node delays for a random number of backoff periods in  $[0, 2^{m_0} - 1]$ , where  $m_0$  is the **initial backoff exponent**.
- If clear channel assessments (CCA) is idle, the node starts the transmission and waits for an ACK.
- If the channel is busy, the procedure is repeated increasing the backoff window until a **maximum backoff exponent  $m_b$** .
- After a **maximum number of backoffs  $m$**  the packet is discarded.
- In case of collision the procedure is restarted and repeated until a **retry limit  $n$**





# Markov Chain Model



- Three dimensional Markov chain  $(s, c, r)$ 
  - $s$ : backoff stage
  - $c$ : state of backoff counter
  - $r$ : state of retransmission counter
- Model characteristic parameter
  - $\lambda = 1/q_0$ : traffic condition
  - $m_0, m, m_b, n$ : MAC parameters
- Computed model parameters
  - $\alpha_i$ : busy channel probability during CCA
  - $P_{coll,i}$ : collision probability





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# Model Equations

- Channel access probability

$$\tau_l = \left( \frac{1 - \alpha_l^{m+1}}{1 - \alpha_l} \right) \left( \frac{1 - y_l^{n+1}}{1 - y_l} \right) b_{0,0,0}^{(l)}$$

$$b_{0,0,0}^{(l)} = \begin{cases} \left[ \frac{1}{2} \left( \frac{1 - (2\alpha_l)^{m+1}}{1 - 2\alpha_l} W_0 + \frac{1 - \alpha_l^{m+1}}{1 - \alpha_l} \right) \frac{1 - y_l^{n+1}}{1 - y_l} + (L_s(1 - P_{\text{coll},l}) + L_c P_{\text{coll},l})(1 - \alpha_l^{m+1}) \frac{1 - y_l^{n+1}}{1 - y_l} \right. \\ \left. + \frac{1 - q_l}{q_l} \left( \frac{\alpha_l^{m+1}(1 - y_l^{n+1})}{1 - y_l} + P_{\text{coll},l}(1 - \alpha_l^{m+1})y_l^n + (1 - P_{\text{coll},l}) \frac{(1 - \alpha_l^{m+1})(1 - y_l^{n+1})}{1 - y_l} \right) \right]^{-1}, \\ \text{if } m \leq \bar{m} = m_b - m_0 \\ \\ \left[ \frac{1}{2} \left( \frac{1 - (2\alpha_l)^{\bar{m}+1}}{1 - 2\alpha_l} W_0 + \frac{1 - \alpha_l^{\bar{m}+1}}{1 - \alpha_l} + (2^{m_b} + 1)\alpha_l^{\bar{m}+1} \frac{1 - \alpha_l^{m - \bar{m}}}{1 - \alpha_l} \right) \frac{1 - y_l^{n+1}}{1 - y_l} \right. \\ \left. + (L_s(1 - P_{\text{coll},l}) + L_c P_{\text{coll},l})(1 - \alpha_l^{m+1}) \frac{1 - y_l^{n+1}}{1 - y_l} \right. \\ \left. + \frac{1 - q_l}{q_l} \left( \frac{\alpha_l^{m+1}(1 - y_l^{n+1})}{1 - y_l} + P_{\text{coll},l}(1 - \alpha_l^{m+1})y_l^n + (1 - P_{\text{coll},l}) \frac{(1 - \alpha_l^{m+1})(1 - y_l^{n+1})}{1 - y_l} \right) \right]^{-1}, \\ \text{otherwise,} \end{cases}$$



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# Model Equations

- Busy channel probability

$$\alpha_{\text{pkt},l} = L \sum_{i=0}^{|\Omega_l|-2} \sum_{j=1}^{C_{l,i}} \prod_{k=1}^{i+1} \tau_{k_j} \left( 1 - \prod_{k=1}^{i+1} \alpha_{k_j} \right) \prod_{h=i+2}^{|\Omega_l|} (1 - \tau_{h_j})$$

$$\alpha_{\text{ack},l} = L_{\text{ack}} \sum_{h \in \Omega_0, h \neq l} q_h R_h$$

- Collision probability

$$P_{\text{coll},l} = \Pr[\mathcal{A}_l] + (1 - \Pr[\mathcal{A}_l]) \Pr[\mathcal{B}_l]$$

$$\Pr[\mathcal{A}_l] = 1 - \prod_{k \in \Omega_0, k \neq l} (1 - \tau_k)$$

$$\Pr[\mathcal{B}_l] = 2L \sum_{i=0}^{|\Omega_0 \setminus l|-2} \sum_{j=1}^{C_{l,i}} \prod_{k=1}^{i+1} \tau_{k_j} \left( 1 - \prod_{k=1}^{i+1} \alpha_{k_j} \right) \prod_{h=i+2}^{|\Omega_l|} (1 - \tau_{h_j})$$



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# Routing Model

- Let  $\lambda \in \mathbb{R}^N$  be a vector of node packet generation rates.
- Let  $\pi_{i,j}$  be the metric associated to link  $l=(i,j)$ .
- The effect of the routing is described by a matrix  $T$  such that:

$$T_{i,j} = \Pr \left[ \pi_{i,j} = \max_{V_h \in \Gamma_i} \pi_{i,h} \right]$$

$$\tilde{T}_{i,j} = T_{i,j} R_l$$

- In stationary conditions, the vector of traffic rates can be described by a system of flow balance equations:

$$Q = Q \tilde{\mathbf{T}} + \lambda \quad Q \in \mathbb{R}^N$$





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# Performance Indicators

- Reliability
  - Homogeneous link quality

$$R_l = 1 - (P_{cf,l} + P_{cr,l})$$

$$P_{cf,l} = \frac{\alpha_l^{m+1} (1 - (P_{coll,l}(1 - \alpha_l^{m+1}))^{n+1})}{1 - P_{coll,l}(1 - \alpha_l^{m+1})}$$

$$P_{cr,l} = (P_{coll,l}(1 - \alpha_l^{m+1}))^{n+1}$$

- Delay
  - Measured for successfully received packets

$$\mathbb{E}[D_l] = \sum_{j=0}^n \Pr(\mathcal{C}_j | \mathcal{C}) \mathbb{E}[D_{l,j}]$$

$$\Pr(\mathcal{C}_j | \mathcal{C}) = \frac{(1 - P_{coll,l}(1 - \alpha_l^{m+1})) P_{coll,l}^j (1 - \alpha_l^{m+1})^j}{1 - (P_{coll,l}(1 - \alpha_l^{m+1}))^{n+1}}$$

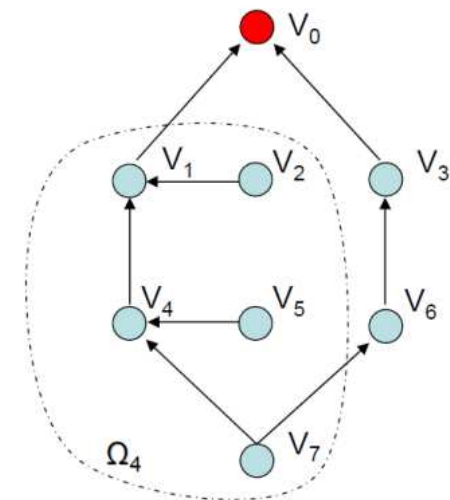
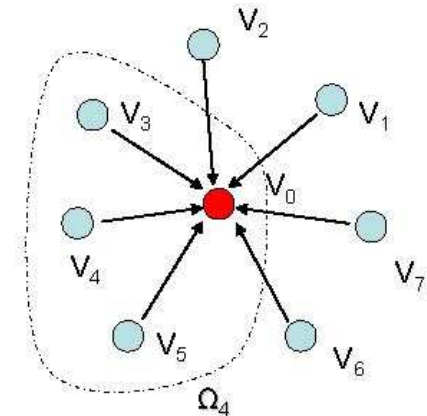




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# Simulation Setup

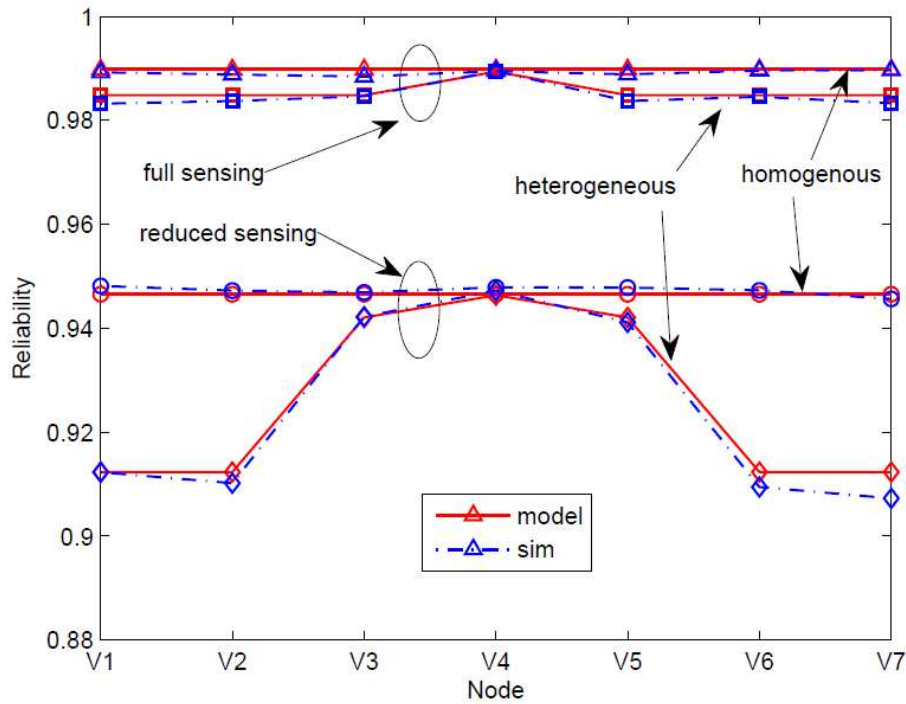
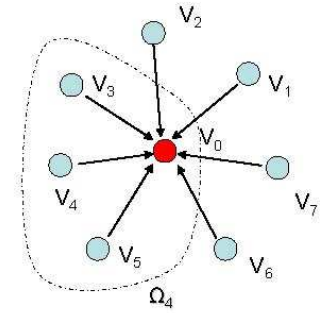
- Single-hop Analysis
  - Traffic condition
    - Homogeneous ( $\lambda_i=5$  pkt/s  $i=1,\dots,N$ )
    - Heterogeneous ( $\lambda_4=20$  pkt/s;  $\lambda_i=5$  pkt/s)
  - Sensing range
    - Full sensing ( $|\Omega_i|=N$   $i=1,\dots,N$ )
    - Reduced sensing ( $|\Omega_i|=3$   $i=1,\dots,N$ )
- Multi-hop Analysis
  - Path selection
    - Path 1 ( $V_7-V_4-V_1-V_0$ )
    - Path 2 ( $V_7-V_6-V_3-V_0$ )
  - Coupling
    - Coupled paths
    - Isolated paths



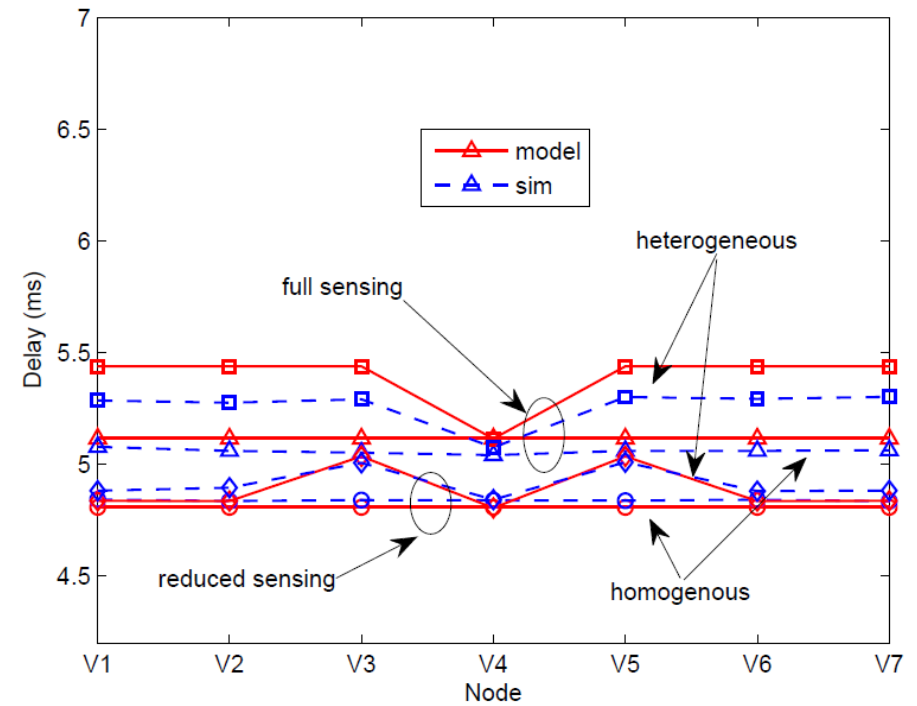


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# Single-hop Performance



Reliability increases when the number of contenders and hidden terminals decreases



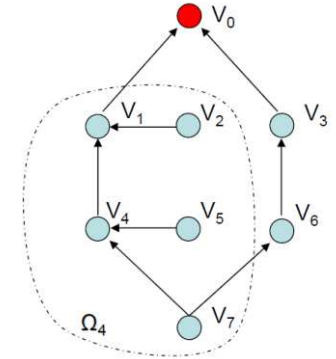
Delay decreases when the number of contenders decreases



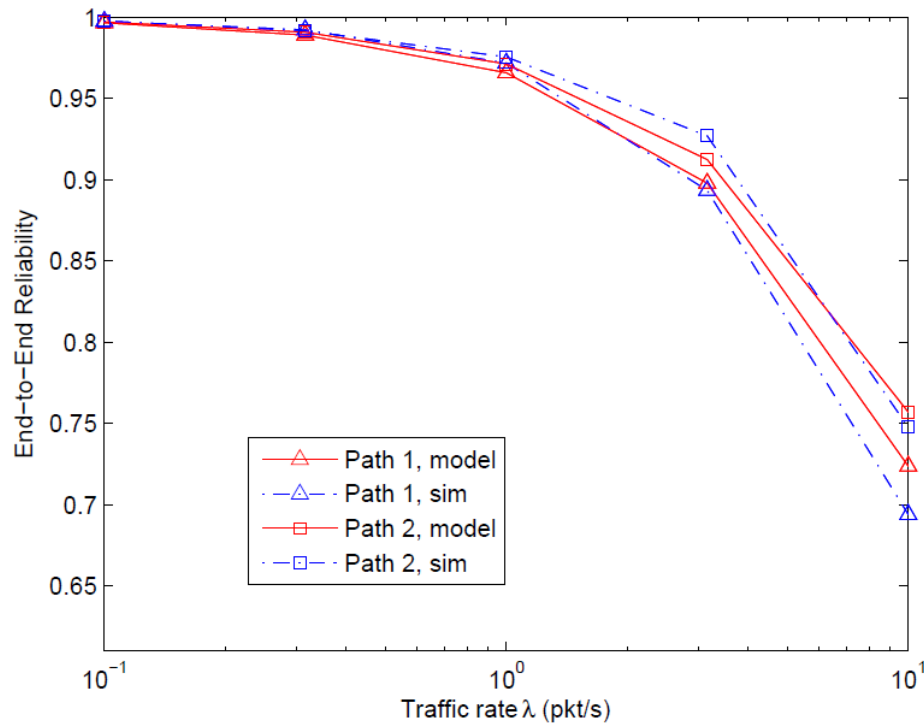


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# Reliability Analysis

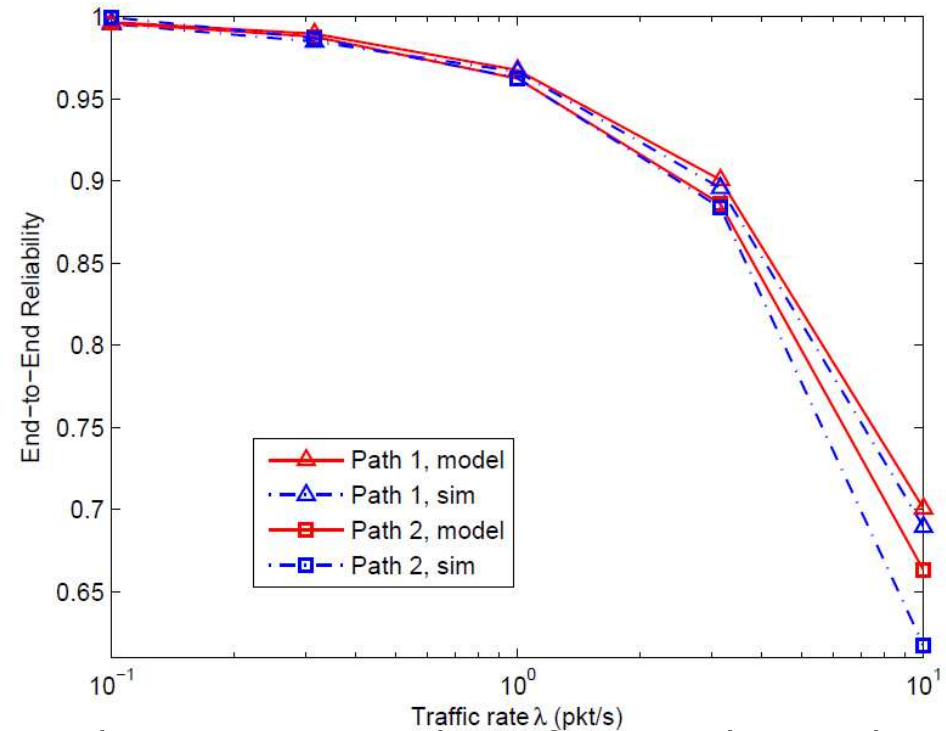


Isolated paths



Number of contenders in each time unit is lower in Path 2  $\rightarrow$  better performance

Coupled paths



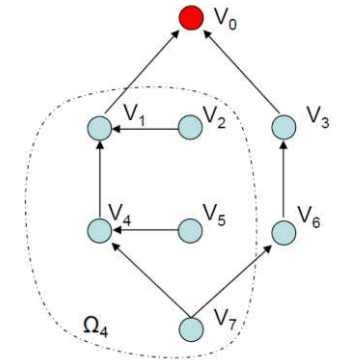
The average number of contenders reduces by forwarding packets to Path 1



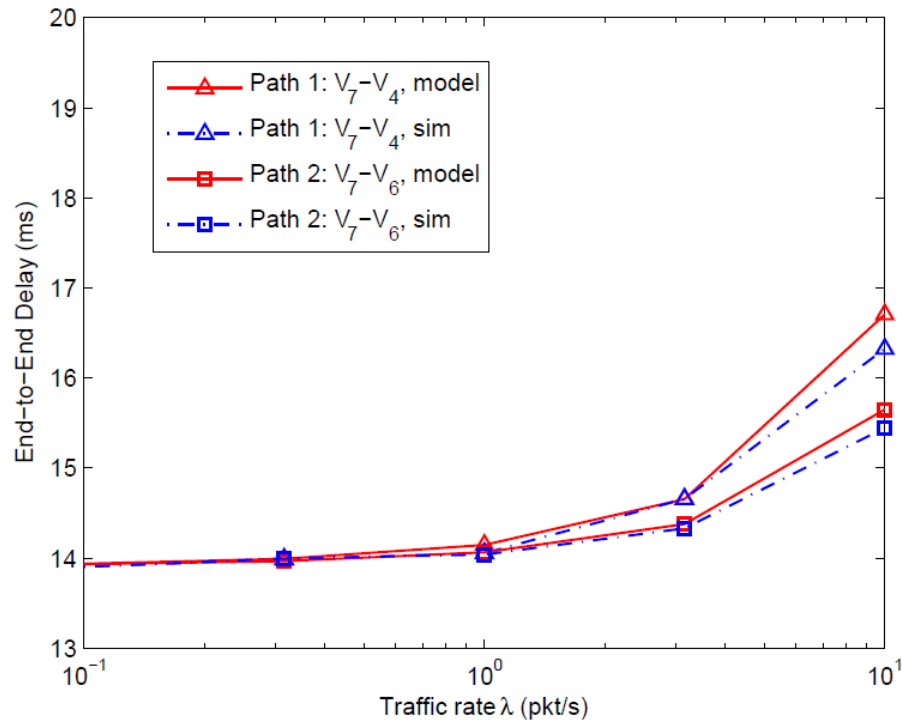


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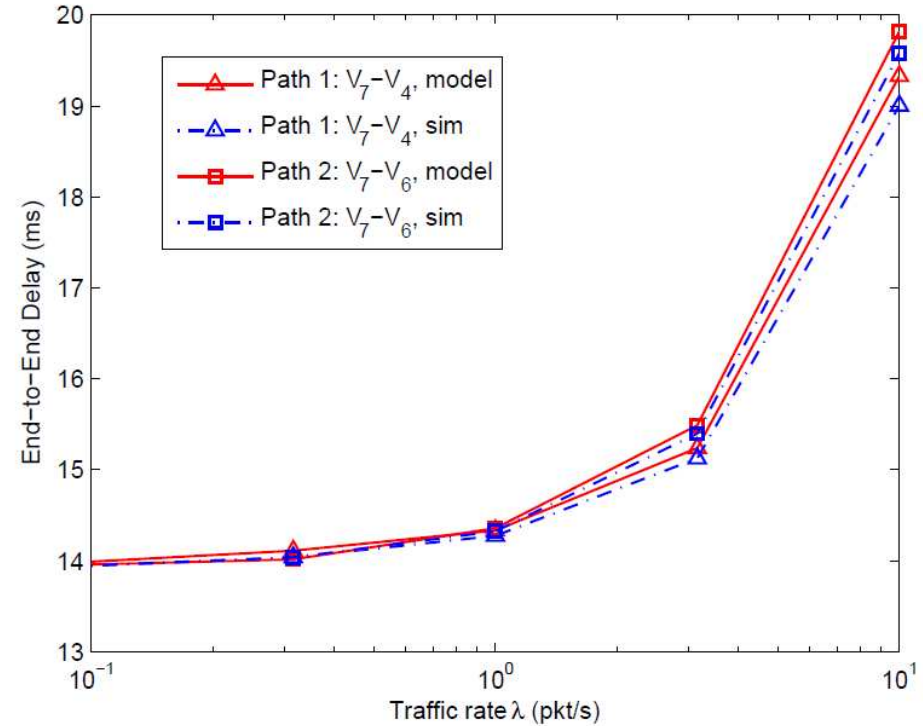
# Delay Analysis



Isolated paths



Coupled paths





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# Considerations

- Reliability and delay at MAC layer impact on the performance of routing algorithms
- Metrics based on reliability and delay may lead to unbalanced traffic distribution, by forcing the the forwarded traffic to dominant paths.
- Open problem:
  - How to prove the unicity of the solution of model equations





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# Thanks for the attention

Questions?

## Reference:

- P. Di Marco, P. Park, C. Fischione, and K. H. Johansson: *Analytical Modelling of IEEE 802.15.4 for Multi-hop Networks with Heterogeneous Traffic and Hidden Terminals*. IEEE GLOBECOM 2010.

